

A Research Councils UK High End Computing Service

Parallelisation of CABARET

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Outline

- Project background
- What is CABARET?
- Applications of CABARET
- Why develop the CABARET code?
- Code description and performance
- Conclusion





- Started March 2009 working 50%
- Collaboration with the Whittle Laboratory, University of Cambridge Department of Engineering
- First dCSE project requiring full parallelisation of a code
- Ends Feb 2011





 Compact Accurately Boundary Adjusting high-Resolution Technique (CABARET)

- Method is based on an extension of the original second-order Upwind Leapfrog three-time-level advection scheme (Roe, 1993)
- Two-time-level non-oscillatory scheme for quasilinear hyperbolic conservation laws





Applicable to compressible unsteady Navier-Stokes
 Flow problems ranging from acoustic wave propagation
 and vortical flows to shock wave interaction

• In Re~10000 calculations the method gives a very good convergence without additional preconditioning, down to Mach numbers as low as M~0.05-0.1





CABARET LES ocean modelling

• Because of hydrodynamic instability, the boundary layer, which occurs at the left (western) domain boundary, is separated and a free jet in the eastward direction is developed

• Time-averaged stream function in the top layer of the 3-layer model, CABARET is \sim 30 times more efficient!



3-D backward facing step

Test case Low Mach number turbulent flow around a 3-D backward facing step



Why develop parallel CABARET?

- Original Fortran 90 code was developed around 1998
- Can handle up to a million grid points
- To resolve higher resolution grids required in high fidelity LES simulations where Re>100000 we require 10 million grid points
- Locality of algorithm will lend itself well to distributed processing – develop parallel code





Data decomposition

CABARET is similar to a finite difference / finite volume calculation

 Parallel version of the code already exists for a structured orthogonal grid

- This project is concerned with developing a parallel version based on an irregular hexahedral grid
- Next step is implementing a tetrahedral cell structure





Parallel Data Decomposition

- For the data decomposition we use calls to the graph partitioner Metis
- Metis produces minimal edge cuts for each partition
- thus minimising MPI communications





E.g. Four part decomposition for the backstep case



- PHASE1 conservative predictor step
- VISCOSITY computation of the cell centred viscous terms
- PHASE2 / MODULE- extrapolation step where the local cellbased characteristic splitting is performed
- BOUND applying physical boundary conditions for the boundary cells
- PHASE3 conducting the conservative corrector step







- All calculations are local to a cell and it's six nearest neighbour cells
- All main loops involve NSIDE calculations apart from phase1
- Finite volume calculation in phase1 loops over the APEXes



 Irregular decomposition – local numbering on each cpu is non-contiguous

- Global to local mapping for SIDEs, APEXes, CELLs and boundary SIDEs
- Decomposition is optimised for the SIDEs





Partition connectivity

• Each cpu stores connectivity between neighbouring cpus and their SIDEs - NEIGH(I)

Connectivity local to each cpu –
 cpu 1 : connect(2:3)=1, cpu 2 : connect(1)=1, cpu 3 : connect(1)=1







SIDE connectivity

- K=SIDELINK(L,I) where K and L are local SIDE numbers and I is the neighbouring partition
- K and L both map to the same global SIDE number







Vectorisation

All the main loops are SIDE based
 DO L=1,NSIDE

```
IF (GEMSIDECELL(L,1)/=0)...
IF (GEMSIDECELL(L,3)==1234)...
END DO
```

```
will not vectorise !
```





Pointers

- Main arrays for the grid data, flux-type and conservative terms are allocatable
 - REAL(KIND=8), TARGET, ALLOCATABLE :: CELL(:)
 - INTEGER, TARGET, ALLOCATABLE :: GEMCELLSIDE(:)

Use Fortran90 data types

- TYPE TRANSFERS
- INTEGER, POINTER :: INTBLOCK(:)
- REAL(KIND=DP), POINTER :: REALBLOCK(:)
- END TYPE TRANSFERS
- TYPE (TRANSFERS) :: TRANSFER1
- ALLOCATE(TRANSFER1%INTBLOCK(4*NCELL))
- ALLOCATE(TRANSFER1%REALBLOCK(4*NCELL))





Transfers with pointers

- Set pointers
 - TRANSFER1%INTBLOCK=>GEMCELLSIDE
 - TRANSFER1%REALBLOCK=>CELL
- Can pass with separate calls
 - CALL MPI_ISSEND(TRANSFER1%INTBLOCK,...
 - CALL MPI_ISSEND(TRANSFER1%REALBLOCK,...
- Why not send the TRANSFER data type as an MPI_TYPE
 - CALL MPI_TYPE_CREATE_STRUCT((2, BLOCKCOUNTS, OFFSETS, OLDTYPES,TRANSFERTYPE, IERR)
 - CALL MPI_TYPE_COMMIT(TRANSFERTYPE, IERR)





Initial performance

• Test initial parallelisation with the following HECToR compilers and optimisation flags

- Cray FFLAGS = -O3 -Oaggress -Omsgs
- Pgf90 FFLAGS = -Minfo -Mneginfo -Mextend -fast -Munroll=n:4 -Mipa=fast,inline
 -O3 -tp barcelona-64
- Pathscale FFLAGS = -Ofast -LNO:full_unroll=4 -march=barcelona
 -OPT:malloc_algorithm=1 -LNO:simd_verbose=ON
- Gfortran FFLAGS = -march=barcelona -ffast-math -funroll-loops -O3 -ffixed-linelength-72 -ftree-vectorizer-verbose=2
- See http://www.hector.ac.uk/cse/reports/compilers.php for compiler performance results on a variety of other codes





What takes most time?

One iteration of CABARET algorithm







Performance





Backward facing step case (fixed problem size) with NAPEX=111741 NCELL =100000 NSIDE =311400 270 iterations



Performance





Backward facing step case (fixed problem size) with NAPEX=111741 NCELL =100000 NSIDE =311400 270 iterations



3-D backward facing step

Paraview plot of x component of velocity





Backward facing step case with NAPEX=111741 NCELL =100000 NSIDE =311400 mature solution



Conclusion

- CABARET method up to 30 times more efficient for some CFD applications on coarser grids
- For Re~10000 and subsonic M method gives good convergence without preconditioning
- Parallel code uses irregular domain decomposition
- Main loops will not vectorise further work!



