

Performance Analysis: Single-Node

The importance of profiling

- **The Golden Rules of profiling:**
 - **Profile your code**
 - The compiler/runtime will **NOT** do all the optimisation for you.
 - **Profile your code yourself**
 - Don't believe what anyone tells you. They're wrong.
 - **Profile on the hardware you want to run on**
 - Don't profile on your laptop if you plan to run on HECToR.
 - **Profile your code running the full-sized problem**
 - The profile will almost certainly be qualitatively different for a test case.
 - **Keep profiling your code as you optimise**
 - Concentrate your efforts on the thing that slows your code down.
 - This will change as you optimise.
 - So keep on profiling.

What we talk about here

- **You've already seen how to profile codes**
- **Here we will show you some extra features**
 - These are most useful for improving single node performance, e.g.
 - a serial code, running on just one core
 - a single rank of a parallel (prob. MPI) job
 - that rank is running on one core of a node
 - an OpenMP code running threads on the cores of a single node
 - one rank of a hybrid MPI/OpenMP code
 - that rank is running threads on the cores of a single node
 - So we are not looking at profiling usage of the network or I/O here.
- **What we are looking at:**
 - loop-level profiling
 - hardware performance counters
 - hybrid MPI/OpenMP jobs
 - memory footprint and overheads

Loop statistics

- **Function-level tracing**
 - Identifies routines taking most time; candidates for optimisation
 - But, where in the routine is time spent?
 - on one single task or spread across many small tasks?
- **Loop-level tracing gives finer granularity**
 - Most HPC codes are based round loopnests
 - If one loopnest dominates runtime, can think about rewriting it
 - If lots of loopnests contribute, maybe look at combining them
- **Information provided**
 - Inclusive loop times
 - How much time was spent executing that loopnest
 - Loop tripcounts (min, mean, max)
 - Can be used to give compiler hints to allow additional optimisation
 - tripcounts very interesting if considering a GPU port
- **Only available with Cray compiler**
 - Compiler adds some additional features into binary executable.
 - Should be done as separate experiment
 - compiler optimizations are restricted with this feature
 - this will affect runtime and probably also skew the profile

Collecting loop statistics

- **Get the right modules**
 - `PrgEnv-cray` because this is a CCE-specific feature
 - `perftools` because we need profiling
- **Re-compile & link code**
 - Add `-h profile_generate` flag to compile and link lines
- **Instrument binary executable for tracing using one of:**
 - `pat_build -u my_program`
 - traces all user functions; fine for small codes
 - `pat_build -w -T func1_,func2_... my_program`
 - traces specified functions, better for larger programs
 - watch out for trailing underscores (`pat_build` will say how many functions traced)
 - `pat_build` using an existing `apa` file.
- **Run the instrumented application**
 - `aprun <options> my_program+pat`
- **Create report with loop statistics**
 - `pat_report my_program+pat+<etc>.xf > loops_report`

Default report

Table 2: Loop Stats by Function (from `-hprofile_generate`)

Loop Incl Time Total	Loop Hit	Loop Trips Avg	Loop Trips Min	Loop Trips Max	Function=/.LOOP[.] PE=HIDE
6.067347	100	25	0	25	sweepy_.LOOP.1.li.32
6.067182	2500	25	0	25	sweepy_.LOOP.2.li.33
5.877503	50	25	0	25	sweepz_.LOOP.05.li.48
5.877408	1250	25	0	25	sweepz_.LOOP.06.li.49
2.971128	50	25	0	25	sweepx2_.LOOP.1.li.28
2.971051	1250	25	0	25	sweepx2_.LOOP.2.li.29
2.947443	50	25	0	25	sweepx1_.LOOP.1.li.28
2.947376	1250	25	0	25	sweepx1_.LOOP.2.li.29
2.247618	187500	107	0	107	riemann_.LOOP.2.li.63

Perfectly nested loops:

- Incl time almost same
- Outer Hit*Avg = Inner Hit
- Consecutive source lines

Hardware performance counters

- **CrayPAT can interface with Cray XE6's HWPCs**
 - Gives extra information on how hardware is behaving
 - Very useful for understanding (& optimising) application performance
- **Provides information on**
 - hardware features, e.g. caches, vectorisation and memory bandwidth
- **Available on per-program and per-function basis**
 - Per-function information only available through tracing
- **Number of simultaneous counters limited by hardware**
 - 4 counters available with AMD Opteron Interlagos
 - If you need more, you'll need multiple runs
- **Most counters accessed through the PAPI interface**
 - Either native counters or derived metrics constructed from these

Hardware counters selection

- **HWPCs collected using CrayPAT**
 - Compile and instrument code for profiling as before
- **Set `PAT_RT_HWPC` environment variable at runtime**
 - e.g. in the your job script
 - Hardware counter events are **not** collected by default (except with APA)
- **`export PAT_RT_HWPC=...`**
 - either a list of named PAPI counters
 - or `<set number>` = a pre-defined (and useful) set of counters
 - recommended way to use HWPCs
 - there are around 20 groups to choose from
 - To see them:
 - `pat_help -> counters -> amd_fam15h -> groups`
 - `man hwpc`
 - `more ${CRAYPAT_ROOT}/share/CounterGroups.amd_fam15h`
- **HW counter information shown by `pat_report`:**
 - Raw data
 - Derived metrics
 - Desirable thresholds

Technical term for Interlagos

Predefined Interlagos HW Counter Groups

Default is number 23 with CrayPAT APA procedure

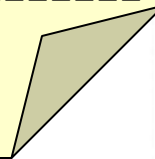
0: Summary, instructions metrics
1: Summary with TLB metrics
2: L1 and L2 Metrics
3: Bandwidth information
5: Floating operations dispatched
6: Cycles stalled, resources idle
7: Cycles stalled, resources full
8: Instructions and branches
9: Instruction cache

11: Floating point ops dispatched
12: Dual pipe flops dispatched
13: Floating point operations SP
14: Floating point operations DP
19: Prefetchs
20: FP, D1, TLB, MIPS
21: FP, D1, TLB, Stalls
22: D1, TLB, MemBW
23: FP, D1, D2, and TLD

Example: HW counter data and derived metrics



```
=====
USER / fftstp_
-----
Time%                                4.9%
Time                                3.441629 secs
Imb. Time                            1.170746 secs
Imb. Time%                            26.2%
Calls                                0.117M/sec      412665.0 calls
...
HW FP Ops / User time                3016.556M/sec   10670175126 ops   29.0%peak(DP)
Total DP ops                         3016.556M/sec   10670175126 ops
Computational intensity                1.16 ops/cycle   1.04 ops/ref
MFLOPS (aggregate)                   96529.78M/sec
TLB utilization                        5420.38 refs/miss 10.587 avg uses
D1 cache hit,miss ratios               94.8% hits       5.2% misses
D1 cache utilization (misses)          19.30 refs/miss  2.412 avg hits
D2 cache hit,miss ratio                95.4% hits       4.6% misses
D1+D2 cache hit,miss ratio             99.8% hits       0.2% misses
D1+D2 cache utilization                415.17 refs/miss 51.897 avg hits
D2 to D1 bandwidth                    5201.921MB/sec   19294072056 bytes
Average Time per Call                  0.000008 secs
CrayPat Overhead : Time                55.7%
=====
```



Some hints on interpreting the data

● TLB utilization

- Memory loaded in pages: 4kB standard (could use larger hugepages)
- e.g. 512 x 8-byte double precision floats
- So if every double was used once, expect 512 refs/miss
 - Less than 512 shows poor use; more than 512 is good (5420.38 excellent)
 - N.B. $\langle \text{avg uses} \rangle = \langle \text{refs/miss} \rangle / 512$

● D1 cache utilization

- Level 1 cache line is 64 contiguous bytes, e.g. 8 x 8-byte doubles
- So if every double was used once, expect 8 refs/miss
 - Corresponds to hit ratio of 87.5% [i.e. $100 * (1 - 1 / \langle \text{refs/miss} \rangle)$]
 - N.B. $\langle \text{avg uses} \rangle = \langle \text{refs/miss} \rangle / 8$
 - Less than 8 (or 87.5%) shows poor use
 - Rule of thumb: want this to be 20 (or 95%) or more

● D1+D2 cache hit ratio

- Should be high (rule of thumb is more than 97%);

CrayPAT observations and suggestions

D1 + D2 cache utilization: 39.8% of total execution time was spent in 4 functions with combined D1 and D2 cache hit ratios below the desirable minimum of **97.0%**. Cache utilization might be improved by modifying the alignment or stride of references to data arrays in these functions.

D1_D2_cache_hit_ratio	Time%	Function
56.8%	12.0%	calc3_
77.9%	6.4%	calc2_
95.7%	1.4%	calc1_
96.3%	20.0%	<u>calc3 .LOOP@li.80</u>

TLB utilization: 19.6% of total execution time was spent in 3 functions with fewer than the desirable minimum of **512** data references per TLB miss. TLB utilization might be improved by modifying the alignment or stride of references to data arrays in these functions.

LS_per_TLB_DM	Time%	Function
2.56	12.0%	calc3_
5.32	6.3%	calc2_

Interpreting the performance numbers

- **Performance numbers are an average over all ranks**
 - explains non-integer values
- **This does not always make sense**
 - e.g. if ranks are not all doing the same thing:
 - Master-slave schemes
 - MPMD apruns combining multiple, different programs
- **Want them to only process data for certain ranks**
 - `pat_report -sfilter_input='condition' ...`
 - `condition` should be an expression involving `pe`, e.g.
 - `pe<1024` for the first 1024 ranks only
 - `pe%2==0` for every second rank
 - This option is also useful for large `.ap2` or `.xf` data files
 - Compiling a report may take too long
 - or exceed the available memory

OpenMP data collection and reporting

- **Give finer-grained profiling of threaded routines**
 - Measure overhead incurred entering and leaving
 - Parallel regions
 - `#pragma omp parallel`
 - Work-sharing constructs within parallel regions
 - `#pragma omp for`
- **Timings and other data now shown per-thread**
 - rather than per-rank
- **OpenMP tracing enabled with `pat_build -gomp ...`**
 - CCE, PGI: insert tracing points around parallel regions automatically
 - Intel, Gnu, Pathscale: need to use CrayPAT API manually
- **Can also use sampling without API**
 - Get a per-rank view; no per-thread counters
 - Run with `OMP_NUM_THREADS=1` during sampling

OpenMP data collection and reporting

- **Load imbalance for hybrid MPI/OpenMP programs**

- now calculated across all threads in all ranks
- imbalances for MPI and OpenMP combined
 - Can choose to see imbalance in each programming model separately
 - See next slide for details

- **Data displayed by default in pat_report**

- no additional options needed
- Report focuses on where program is spending its time
- Assumes all requested resources should be used
 - you may have reasons not to want to do this, of course

Imbalance options for data display (pat_report -0 ...)

- **These options control how load balance is displayed:**
- **profile_pe_th (default view)**
 - Imbalance based on the set of all threads in the program
 - i.e. imbalance from OpenMP and MPI combined
 - this is best measure to understand code performance
- **profile_pe.th**
 - Highlights imbalance across MPI ranks
 - Thread data for each rank is aggregated
 - max used rather than mean, to highlight under-performers
 - Aggregated thread data merged into MPI rank data
- **profile_th_pe**
 - For each thread, show imbalance over MPI ranks
 - Example: Load imbalance shows where thread 4 in each MPI rank didn't get much work

Memory usage

- **Knowing how much memory each rank uses is important:**
 - What is the minimum number of cores I can run this problem on?
 - given there is 32GB (~30GB usable) of memory per node (32 cores)
 - Does memory usage scale well in the application?
 - Is memory usage balanced across the ranks in the application?
 - Is my application spending too much time allocating and freeing?
- **Profile heap usage using CrayPAT**
 - `pat_build -gheap ...`

Heap statistics

Memory per rank
~30GB usable memory per node

Too many allocs/frees?
Would show up as ETC
time in CrayPAT report

Notes for table 5:

Table option:

-O heap_hiwater

Options implied by table option:

-d am@,ub,ta,ua,tf,nf,ac,ab -b pe=[mmm]

This table shows only lines with Tracked Heap HiWater MBytes > 0.

Table 5: Heap Stats during Main Program

Tracked Heap HiWater MBytes	Total Allocs	Total Frees	Tracked Objects Not Freed	Tracked MBytes Not Freed	PE [mmm]
9.794	915	910	4	1.011	Total
9.943	1170	1103	68	1.046	pe.0
9.909	715	712	3	1.010	pe.22
9.446	1278	1275	3	1.010	pe.43

Memory leaks
Not usually a problem in HPC

Summary

- **Profiling is essential to identify performance bottlenecks**
 - even at single core level
- **CrayPAT has some very useful extra features**
 - can pinpoint and characterise the hotspot loops (not just routines)
 - hardware performance counters give extra insight into performance
 - well-integrated view of hybrid programming models
 - most commonly MPI/OpenMP
 - also CAF, UPC, SHMEM, pthreads, OpenACC, CUDA
 - information on memory usage
- **And remember the Golden Rules**
 - including the one about not believing what anyone tells you

Are there any questions?

```
pat_record(PAT_STATE_OFF);
```